



A Functional Language for Specifying Business Reports

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Outline



Interprise Resource Planning Systems



2 Reports & Report Functions





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Interprise Resource Planning Systems





What are Enterprise Resource Planning Systems?

ERP systems integrate several software components that are essential for managing a business.

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ERP systems integrate

- Financial Management
- Supply Chain Management
- Manufacturing Resource Planning
- Human Resource Management
- Customer Relationship Management

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What do ERP Systems Look Like?





Issues of Many ERP Implementations

Complexity

- processes are specified in general purpose language
- gap between specification and implementation
- large monolithic system

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Inflexibility

- code is duplicated in order to avoid unexpected side effects
- the use of general purpose languages makes customisation expensive
- the (relational) database determines the way data is stored and accessed



Outline





2 Reports & Report Functions





Process-oriented event-driven transaction systems

compact core system

Process-oriented event-driven transaction systems

compact core system • customisable via DSLs

Process-oriented event-driven transaction systems

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Process-oriented event-driven transaction systems





Process-oriented event-driven transaction systems





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Process-oriented event-driven transaction systems





Event Log
event 1
event 2
event 3
event 4
event 5
event 6
event 7
event 8
event 9
event 10









The central data types

- records
- lists



The central data types

- records: events are records
- lists



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Nominal subtyping



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Nominal subtyping



The Report Language – An Example Function

Example

```
reportNames : [String]
reportNames = [pr.name |
cr : CreateReport \leftarrow events,
pr : PutReport = head [ur |
ur : ReportEvent \leftarrow events,
ur.id \equiv cr.id]
]
```



The Report Language – An Example Function





Nominal subtype relation <:

- User defined subtyping partial order on records
- Fixed subtyping relation on built-in types

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 $\alpha.f:\beta \Rightarrow \alpha \rightarrow \beta$



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E.g. record modifier operator _ { $f_1 = -, \dots, f_n = -$ } has type

$$\alpha.f_1:\alpha_1,\ldots,\alpha.f_n:\alpha_n \Rightarrow \alpha \to \alpha_1 \to \ldots \to \alpha_n \to \alpha$$

Record Field Constraints

What do we gain?

- Field names can be used by different record types.
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Example

fullName : (*a.firstName* : **String**, *a.lastName* : **String**) \Rightarrow *a* \rightarrow **String** *fullName x* = *x.firstName* ++ " " ++ *x.lastName*

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Example

 $fullName : (a.firstName : String, a.lastName : String) \Rightarrow a \rightarrow String fullName x = x.firstName ++ " " ++ x.lastName$

$$\begin{array}{l} \textit{setFullName}: (a.\textit{firstName}: \textbf{String}, a.\textit{lastName}: \textbf{String}) \Rightarrow \\ \textbf{String} \rightarrow a \rightarrow a \\ \textit{setFullName name } x = \textbf{let} (\textit{first,last}) = \textit{decompose name} \\ \textbf{in } x \{\textit{firstName} = \textit{first, lastName} = \textit{last}\} \end{array}$$





















ain't easy





Basic idea: unfolding folds

fold f e (x # xs)



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$$f e (x \# xs) = f$$
 x (fold $f e xs$)



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Limitations

- This works well with single folds.
- For nested folds more powerful equations are needed.



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- For nested folds more powerful equations are needed.
 - commutative operations
 - multisets instead of lists



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Conclusions

The Last Slide

What do we have?

- Simple yet powerful data model for ERP
- Purely functional language for extracting & aggregating complex information
- Highly customisable & flexible
- Incrementalisation of report functions

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What are we planning?

- More powerful incrementalisation transformations
- Possibly restricting the language further
- A better cost model